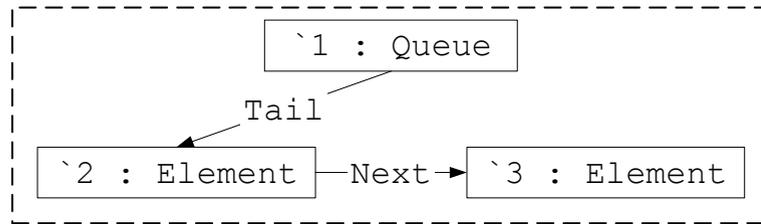


# Solution Sheet 5

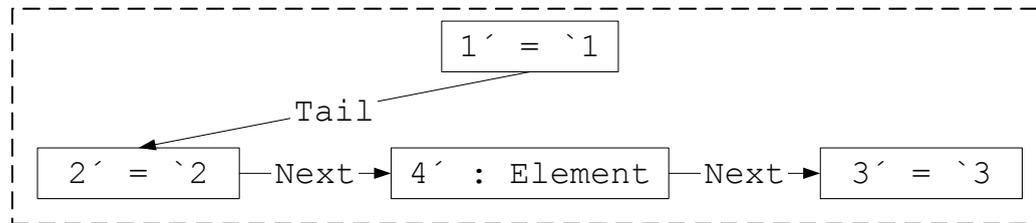
## Exercise 7

## Graph rewrite rule enqueue

production enqueue =



`:=`

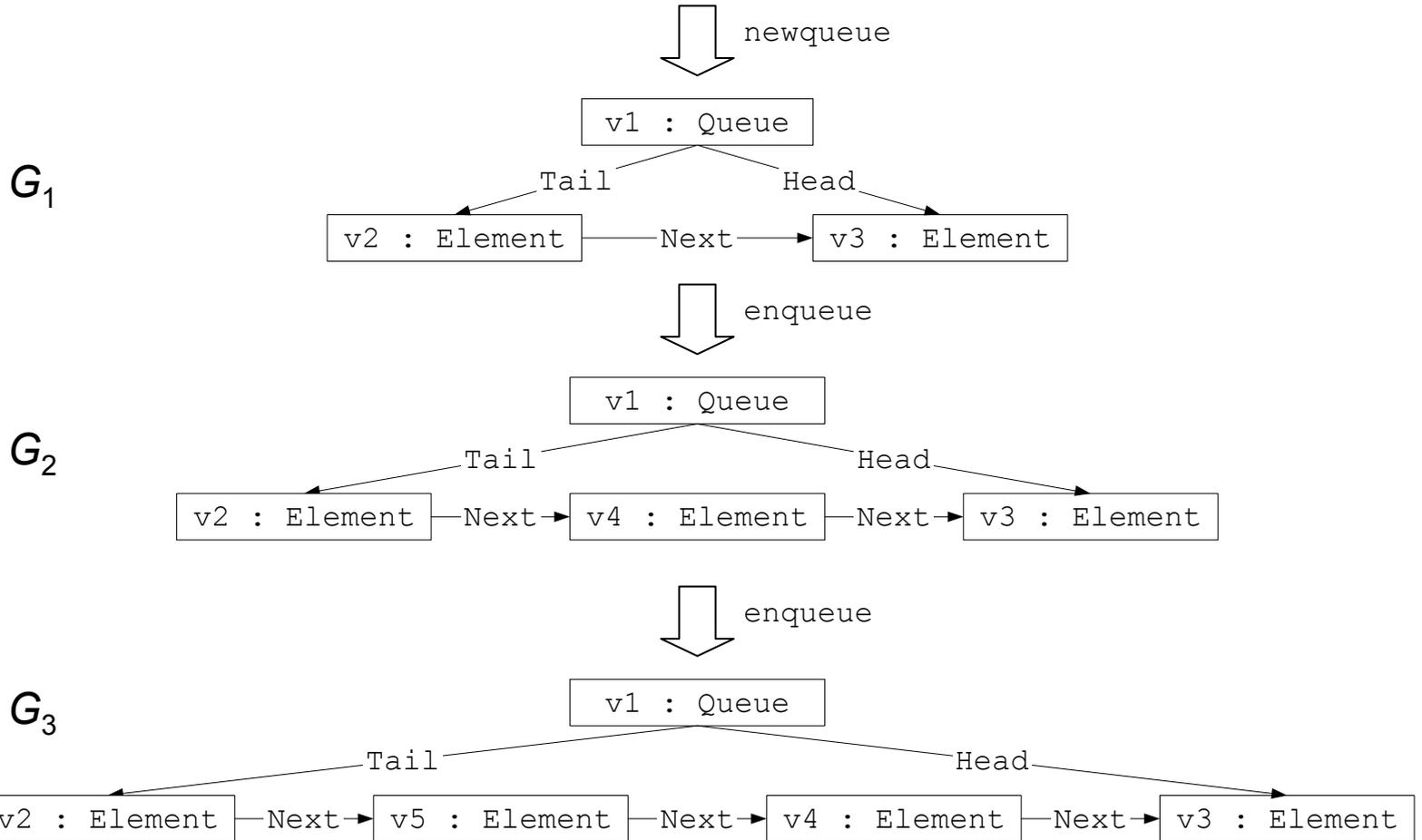


end;

## Formal definition of enqueue

- $r = (L, K, R)$
- $L = (V_L, E_L, I_L)$ 
  - »  $V_L = \{1, 2, 3\}$
  - »  $E_L = \{(1, \text{Tail}, 2), (2, \text{Next}, 3)\}$
  - »  $I_L = \{(1, \text{Queue}), (2, \text{Element}), (3, \text{Element})\}$
- $K = (V_K, E_K, I_K)$ 
  - »  $V_K = \{1, 2, 3\}$
  - »  $E_K = \{(1, \text{Tail}, 2)\}$
  - »  $I_K = \{(1, \text{Queue}), (2, \text{Element}), (3, \text{Element})\}$
- $R = (V_R, E_R, I_R)$ 
  - »  $V_R = \{1, 2, 3, 4\}$
  - »  $E_R = \{(1, \text{Tail}, 2), (2, \text{Next}, 4), (4, \text{Next}, 3)\}$
  - »  $I_R = \{(1, \text{Queue}), (2, \text{Element}), (3, \text{Element}), (4, \text{Element})\}$

# Derivation



## Formal definition of $G_2$

- $G_2 = (V_2, E_2, I_2)$ 
  - »  $V_2 = \{v1, v2, v3, v4\}$
  - »  $E_2 = \{(v1, Tail, v2), (v1, Head, v3), (v2, Next, v4), (v4, Next, v3)\}$
  - »  $I_2 = \{(v1, Queue), (v2, Element), (v3, Element), (v4, Element)\}$

## Application of enqueue to $G_2$

- Definition of an isomorphism  $h : L \rightarrow G_L$ , where  $G_L$  is a partial graph of  $G_2$ 
  - »  $G_L = (V_{GL}, E_{GL}, I_{GL})$ 
    - ⇒  $V_{GL} = \{v1, v2, v4\}$
    - ⇒  $E_{GL} = \{(v1, \text{Tail}, v2), (v2, \text{Next}, v4)\}$
    - ⇒  $I_{GL} = \{(v1, \text{Queue}), (v2, \text{Element}), (v4, \text{Element})\}$
  - »  $h : V_L \rightarrow V_{GL} = \{(1, v1), (2, v2), (3, v4)\}$
  
- Intermediate graph
  - $H = G_2 \setminus (h(L) \setminus h(K)) = G_2 \setminus (\emptyset, \{(v2, \text{Next}, v4)\}, \emptyset)$
  - »  $H$  is obtained from  $G_2$  by deleting the edge  $(v2, \text{Next}, v4)$
  
- Final graph
  - $G_3 = H \oplus h'(R \setminus K) =$
  - $H \oplus h'(\{4\}, \{(2, \text{Next}, 4), (4, \text{Next}, 3)\}, \{(4, \text{Element})\}) =$
  - $H \oplus (\{v5\}, \{(v2, \text{Next}, v5), (v5, \text{Next}, v4)\}, \{(v5, \text{Element})\})$
  - »  $G_3$  is obtained from  $H$  by adding the new node  $v5$  and its incoming and outgoing Next edges

## Other variants of graph rewrite rules

- ❑ Replacement of subgraphs rather than partial graphs
  - ⇒ Two versions of `enqueue` required
    - `enqueue1`: `enqueue` into empty queue (including `Head` edge)
    - `enqueue2`: `enqueue` into non-empty queue (without `Head` edge)
- ❑ Homomorphisms rather than isomorphisms
  - ⇒ Does not extend applicability of graph rewrite rules
    - Only `Element` nodes could be identified
    - In case of identification, a cycle of `Next` edges (of length 1 or 2) would have to be present in the host graph
    - There is no rule which creates a cycle