

6. Preparation for Practice: Translation to Programming Languages

Aims

- ❑ Mapping of architectural concepts to mainstream programming languages
- ❑ Translation can be done by tools as part of an integrated Software Development Environment
 - Architecture → Module Code Frames
 - then using Programming Environment or
 - Architecture Tool – Integrator – Programming Tool
- ❑ Do translation by hand
 - thereby extending architectural knowledge

Contents of Chapter 6

- 6.1. Problems of Translation
- 6.2. Languages with Independent Units (FORTRAN, C)
- 6.3. Block-structured Languages without Modules (Standard Pascal)
- 6.4. Languages with Modules (Ada)
- (6.5. OO Languages (Java))
- 6.6. Summary of Chapter

Level of Programming Language

- High-level languages with rich architectural concepts (e.g. Ada 95):
translation easy, but Programming Language \neq Architectural Language

Low-level (FORTRAN 66, C, Assembler):
translation difficult but possible

- First question: Has programming language module construct ?
i.e. aggregated interface, hidden bodies
else aggregate mentally
corresponding layout and comments
discipline
- Second question: Has language import relations
else simulation as above
- Third question: Has language consistency constraints ?
else discipline

Clarifying Topics

- ❑ Module construct of programming language:
aggregated interface + body
module types f, ado, adt are “semantical“ uses
- ❑ Import construct of programming language
import relation “types“ are “semantical“ uses
- ❑ Translation never impossible
but more or less difficult
- ❑ However: range of mapping depends on language
not possible to map recursion on FORTRAN 66
not possible to map OO on C
So we have to regard “restricted“ architectures

"Classes" of Programming Languages

Regard in this chapter 3 classes:

- ❑ languages without module construct
flat / no hierarchies, arbitrary relations, without checking consistency,
mostly static storage allocation
representatives: FORTRAN 6x, C, Cobol 6x / 7x, Basic, Assembler
- ❑ block-structured languages without module construct
locality by nesting, visibility / scope, no independent compilation
dynamic storage allocation (stack, heap)
representatives: Standard-Pascal, Algol 60, Algol 68 etc.
- ❑ new programming languages with module concepts
independent compilation, more or less checking between units,
multi-paradigmatic
representatives: Ada 83 and 95, Modula-2 / Modula-3, Turbo-Pascal, C++ etc.

Translating to FORTRAN IV

- ❑ FORTRAN IV as representative of languages with independent units, independently compileable
- ❑ Units are not modules but units from which a module body is composed of
- ❑ These languages have subprograms working a global data space (in FTN separable to different COMMOMs)
- ❑ Now: Translation of f, ado, adt modules and module relations locality / general usability

Interface f Module

```

C *****
C * functional module Graph_of_Functions is *
C *   Input data in parameter list *
C *   Output datum is the plotter file *
C * * *
C *   procedure POLYLN (X, Y: inARRAY_RT; *
C *                       XTEXT, YTEXT, UETEXT: in STRING); *
C *   procedure INTLN (...); *
C *   ... *
C *   Semantics of Graph_of_Functions: ... *
C *   ... *
C * end Graph_of_Functions; *
C * ----- *

```

So, interface is just imagination

Body f Module

```

C      module body Graph_of_Functions is -----*
C
C      SUBROUTINE POLYLN (X, Y, XTEXT, YTEXT, UETEXT)
C          REAL X(100), Y(100)
C          INTEGER XTEXT (20), YTEXT(20), UETEXT (20)
C          ...
C          RETURN
C
C      END
C
C      SUBROUTINE INTLN (X, Y, XTEXT, YTEXT, UETEXT)
C          REAL X(100), Y(100)
C          INTEGER XTEXT (20), YTEXT(20), UETEXT (20)
C          ...
C          RETURN
C
C      END
C
C      ...      further "local" subroutines: sealing, plotting axes etc.
C
C      end Graph_of_Functions;
C
C      *****

```

only units of module body are FORTRAN

Interface ado Module

```

C *****
C abstract data object module INTEGER_STACK is
C *   procedure PUSH (X: in INTEGER);
C *   procedure POP;
C *   function READ_TOP return INTEGER: - - - RDTOP
C *   function IS_EMPTY return BOOLEAN; - - - ISEMTY
C *   function IS_FULL return BOOLEAN;   - - - ISFULL
C *   ...
C *   Semantics of operations
C *   ...
C end INTEGER_STACK;
C *-----*
```

Body ado Module (1)

```
C      module body INTEGER_STACK is -----*
```

```
      SUBROUTINE PUSH (ELEMENT)
```

```
          INTEGER STACK (100), POINTR, ELEMNT
```

```
          COMMON / STDATA / STACK, POINTR
```

```
          ...
```

```
          RETURN
```

```
      END
```

```
C
```

```
      SUBROUTINE POP
```

```
          INTEGER STACK (100), POINTR
```

```
          COMMON / STDATA / STACK, POINTR
```

```
          ...
```

```
          RETURN
```

```
      END
```

Body ado Module (2)

C

```
INTEGER FUNCTION RDTOP
    INTEGER STACK (100), POINTR
    COMMON / STDATA / STACK, POINTR
    ...
    RETURN
END
```

C

```
...
BLOCKDATA
    INTEGER STACK (100), POINTR
    COMMON / STDATA / STACK, POINTR
    ...
END
```

C

```
end INTEGER_STACK; *****
```

Mapping Locality

FORTRAN programs are an unstructured heap of units (which compose module bodies)

⇒ no module relation contains, local usability, general usability

⇒ simulation by layout, comments, and discipline

```

C ***** the locally importing module
C * functional module USERINPUT is *
C *      procedure INPUT (...); *
C * * *
C *      The semantics of interface operation is: ... *
C *      ... *
C * * *
C * end USERINPUT; -----*
C * * *
C * module body USERINPUT is -----*
C *      local import from CHECKINPUT *
C *      using CHECKSTRING, CHECKNUM; *
C *      Realization of above operations: *
C *      ... *
C * end USERINPUT; *
C * -----*

```

(not recursive situations !)

The Local Module

```

C ***** *
C functional module CHECKINPUT is *
C * is contained in USERINPUT; *
C * procedure CHECKSTRING (...); *
C * procedure CHECKNUM (...); *
C * ... *
C * The semantics of interface operations: ... *
C * ... *
C end CHECKINPUT; ----- *
C * module body CHECKINPUT is ----- *
C * * *
C ... *
C end CHECKINPUT; *
C * ----- *
```

General Usability & Other Concepts

- ❑ general usability nearer to FORTRAN IV concepts
- ❑ analogous to above
 - general import clause at the interface or
 - body of using component
- ❑ subsystems by two-fold simulation
- ❑ genericity only by macro mechanisms
- ❑ OO not available: cannot easily be simulated
- ❑ recursion not available

Translation to C

- ❑ every module to 2 files
interface: header file
body: implementation file
- ❑ No module concept but rigid help on
program file management
- ❑ simulation (layouting, comments, discipline) similar to FORTRAN
- ❑ subsystems via double simulation
- ❑ genericity only by macro
- ❑ OO not available
- ❑ concentrate on ado mudules, f analogous
adt modules analogous, better with classes from C++

Ado Interface in C

```
// abstract data object module Buffer *****
# ifndef __ INC_BUFFER_H __
# define __ INC_BUFFER_H __

void enqueue (item x);
void deque (item &x);
bool nonempty ();
bool nonfull ();

// item is the type of buffer elements.
// semantics: enqueue adds an element at the backend if buffer is nonfull,
// ...
// -----
```

Ado Body in C

```
# include "Buffer.h"
//general import Virt_array using all
// abstract data object module body Buffer is -----
// realization: circular storage using another ado module Virt_Array
//administration info for the data within Virt_Array container
static int In = 0
static int Out = 0
static int n = 0
static bool nempty = false
static bool nfull = true
// static data scope is restricted to the implementation file

void enqueue (item x) {
    if (n < N) {
        storeVA (In, x); In = (In +1) % N;
        n ++; nfull = (n < N); nempty = true; }
    else { ...}
    ...
// *****
```

The Used Component Ado

```
// abstract data object module Virt_Array *****  
# ifndef __INC_VIRTARRAY_H__  
# define __INC_VIRTARRAY_H__  
  
void storeVA (int loc, item x);  
void readVA (int loc, item &x);  
  
...  
  
// -----
```

Consistency Conditions (Discussion for FORTRAN and C)

- All consistency constraints of chapter 4 by discipline:
 - » for structure and usability relations (graphic level)
 - » for composing export interfaces to a module of a certain kind
 - » for having consistent import interfaces
 - » for consistency export – import
 - » for consistency import – body

- Furthermore as “modules“ are “simulated“
 - » every module is of one above kinds
 - » bodies are hidden:
 - local procedures not used outside
 - encapsulated data structures not directly read or changed
 - local type declarations not used outside

Translation to Standard Pascal

- ❑ Translation similar
 - » aggregate mentally
 - » compose textually, layout, and use comments
 - » apply discipline
- ❑ Thereby translation of a class of programming languages
 - » block-structured
 - » without modules
 - » locality (nesting) the only vehicle to structure programs
 - » other representatives Algol 60, 68, etc.
- ❑ Translation simpler, as Pascal is 10 years younger? No!
- ❑ Translation of module “types” similar,
have some translation experience
⇒ discuss only adt translation

Specific Difficulty: Order of Declarations

anywhere in constant declarations part

```
const STACKMAX = ... ; (* only for internal use in module Stack *)
```

...

anywhere in type declaration part

```
type StackT = (* only for internal use in module Stack *)
```

```
  record
```

```
    SPACE : array [ 1 ... STACKMAX ] of item
```

```
    INDEX : integer
```

```
  end;
```

anywhere in procedure declaration part: interface and body of Stack

```
(*****
```

```
data type module Stack is
```

```
  type StackT is private
```

```
  procedure PUSH (St: in out StackT; EC: in Item);
```

```
  ...
```

```
end Stack; ----- *)
```

```
(* module body Stack is ----- *)
```

```
  procedure PUSH (St: in out StackT; EC: in Item);
```

```
  ...
```

```
  begin ...
```

```
  end;
```

```
  ...
```

```
(* end Stack; *****)
```

Import Translation of Relations to Standard Pascal

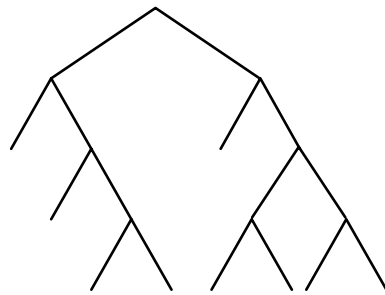
- Locality Translation Idea:
 - contains-relation by nesting
 - local usability by scope / visibility rules } not possible
Pascal has no nesting for modules
- ⇒ locality expressed only by comments

- General Usability:
 - Pascal has no import relations
 - ⇒ general usability by comments

Resulting Program Structure

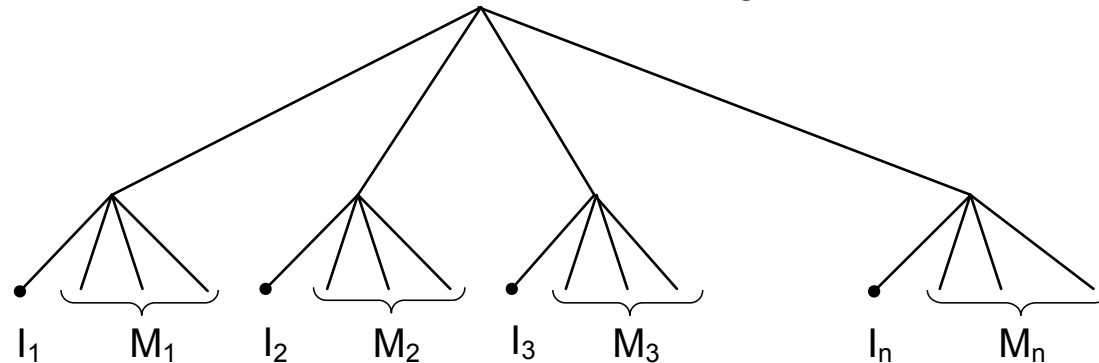
- Where are the “modules“ connected by local and general usability:
In the upmost block
- Structure of a translated program

typical Pascal
program structure

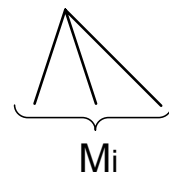


nested
procedures

structure of program
translated from an arch. diagram



I_1 M_1 I_2 M_2 I_3 M_3 I_n M_n



units of a
module body

I_i interface by
comments

+ mutual module relations by comment clause

Example Ada

- ❑ Class of new programming languages with modules, imports, multiparadigmatic language
- ❑ Representatives:
Ada 83, Modula-2, Pascal dialects
- ❑ In addition with OO: C++, Modula-3, Ada 95
- ❑ Language for textual architecture specifications in this lecture for export interfaces, imports
was already Ada 83 or Modula-2 like
- ❑ Discussing translation of module types we again restrict ourselves to adt modules in following example



Adt Module in Ada 83

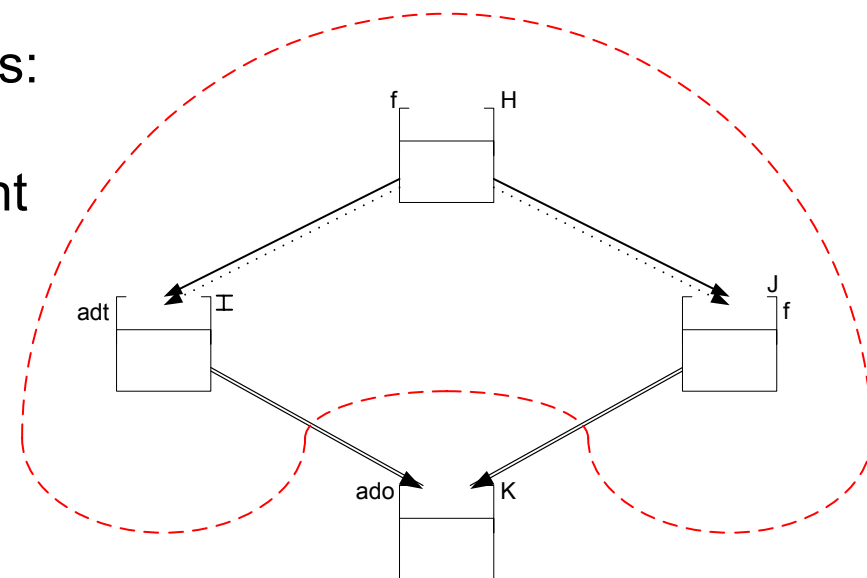
```

-- data type *****
package ITEM_STACK_STENCIL is
  type ITEM_STACK_TYPE is private;
  procedure INITIALIZE (ST: in out ITEM_STACK_TYPE);
  procedure PUSH (EL: in ITEM_TYPE; ST: in out ITEM_STACK_TYPE);
  procedure POP (ST: in out ITEM_STACK_TYPE);
  function READ_TOP (ST: in ITEM_STACK_TYPE) return ITEM_TYPE;
  function IS_EMPTY (ST: in ITEM_STACK_TYPE) return BOOLEAN;
  function IS_FULL (ST: in ITEM_STACK_TYPE) return BOOLEAN;
  ST_UNDERFLOW, ST_OVERFLOW; ST_NOT_INITIALIZED: exception;
  -- semantics: ...
private
  SIZE: constant INTEGER := 100;
  type SPACE_T is array (1..SIZE) of ITEM_TYPE;
  type ITEM_STACK_TYPE is
    record
      SPACE: SPACE_T;
      INDEX: INTEGER range 0..SIZE := 0;
    end record;
end ITEM_STACK_STENCIL; -----
package body ITEM_STACK_STENCIL is -----
  -- implementation of above interface ...
end ITEM_STACK_STENCIL; -- *****

```

Transformation of Locality

- contains relation by nesting:
 - in the declaration part of a module body another one is nested
- big text units result:
 - we have bodies separately: body stubs and separate subunits
(to be developed, compiled)
 - export interface as comments for subunits
 - is – contained clause as comment
- nesting \Rightarrow scope / visibility rules:
 - potential local usability
 - local import clause as comment



The Example 1

```

-- functional *****
package H is
    ...
end H; -----
package body H is -----
    -- local import from I using ...;
    -- local import from J using ...;
    -- abstract data type *****
    package I is
        ...
    end I; -----
    package body I is separate; --*****
    -- functional *****
    package J is
        ...
    end J; -----
    package body J is separate; --*****
    ...
end H; --*****
-----
separate (H);*****
-- abstract data type package I is
--   is contained in H;
--   ...
-- end I; -----
package body I is -----
    ...
end I; --*****
-----
--analogously for J

```

Translation of General Usability

- ❑ general components: library units
- ❑ general usability: Ada import clauses
- ❑ import only of specific resources: comment

The Example 2

```

-- abstract data object *****
package K is
    ...
end K; -----
package body K is -----
    ...
end K; --*****
-----
-- H as above
-----
with K; use K; separate (H) --*****
-- abstract data type package I is
--   is contained in H;
--   ...
--   export interface as comment
-- end I; -----
package body I is -----
    -- general import from K using ...;
    ...
end I; --*****
-----
with K; use K; separate (H) --*****
-- functional package J is
--   is contained in H;
--   ...
--   export interface as comment
-- end J; -----
package body J is -----
    -- general import from K using ...;
    ...
end J; --*****

```

Ada Translation Contd.

- ❑ consistency rules:
most of them by discipline
- ❑ so Ada 83 \neq Architectural Language
Architectural Language conceptually much richer
- ❑ above discussion assumes:
locality \rightarrow nesting and stubs
- ❑ if locality is mapped on library units:
even more simulation and discipline

Subsystems as Aggregated Packages of Ada 83

```

-- data object subsystem *****
package Entry_Coll is
--
  data type *****
  package Entry is
    type Entry_Den_Type is private;
    procedure Create_and_Init (EI: out Entry_Den_Type);
    ...
    -- semantics: ...
  private
    type Entry_Den_Type is ...
  end Entry; --*****
--
  data object *****
  package Coll is
    procedure Store (EI: in Entry_Den_Type);
    ...
    -- semantics: ...
  end Coll; --*****
end Entry_Coll; -----

package body Entry_Coll is -----
  ...
  package body Entry is --*****
    -- Translation of Entry as usual
    -- Body can be developed separately (here only stub)
    ...
  end Entry; --*****
  package body Coll is *****
    -- Translation of Collection as usual
    -- Body can be developed separately (here only stub)
    ...
  end Coll; --*****
begin
  ...
end Entry_Coll; --*****

```

Subsystems as Public Library Units of Ada 95

- ❑ Subsystems interfaces are aggregated
- ❑ Any extension of the interface \implies all clients have to be recompiled
- ❑ Hierarchical libraries of Ada 95 to handle this problem:
 - extension: no recompilation is necessary if client only uses the old interface
- ❑ Mechanism of public child library units can be used
 - » for extending an interface
 - » especially, for extending inheritance hierarchies:
 - inheritance hierarchy = subsystem

Translation of Genericity

- simple genericity
 - discrete types
 - simple types with operations
 - ⇒ generic clauses of Ada 83
 - generic instantiation of generic units
- } restrictions:
no general arch.
within body of
generic unit
- “complex“ types: main application of genericity
 - complex type with operations ⇒ adt
- ⇒ formal generic packages of Ada 95
define the properties which have to be fulfilled

Translation of OO

- most languages of this class have OO
Modula-3, Ada 95, C++

- in Ada 95 inheritance is orthogonal to modularization
⇒ inheritance hierarchies in the interface of a package:
not good if inheritance changes

- better one type (class) = one package
organization of inheritance structure by public child libraries

What We Have Learned

- ❑ Translation discussion valuable:
 - » can architecture language apply to all relevant progr. languages
 - » have deepened understanding of arch. language
- ❑ Architecture Language \neq modern Programming Language
 - » Types of components
 - » Types of relations
 - » Consistency constraints
- ❑ Have learned about Progr. Language Classes
- ❑ Translation can be mechanized by tools